



## **An Enhanced Cost Effective and Scalable Network Architecture for Data Centers**

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### **Abstract**

Large-scale data centers, comprises tens of thousands of computers, require significant total of bandwidth. The network infrastructure typically adopts a hierarchical tree configuration of routing and switching components, with higher levels featuring increasingly specialized and expensive equipment. Despite the implementation of advanced IP switches and routers, the resulting setups often utilize only 50% of the total bandwidth available at the network edge, while still incurring significant costs. Its modular, scalable



design allows expansion with minimal disruption or cost, supporting big data, AI, and other demanding applications. In this architecture balances efficiency, scalability, and compatibility, making it a future-ready, low-cost choice for data centers modernizing for the digital era and expanding data-driven services.

**Keywords:** Data Center, Network, Network Architecture, Scalable Network Architecture, Data, Cost Effective.

## **Introduction**

The size of the DCNs is exponentially growing mainly due to the tremendous surge in the number of the servers. In 2006, Google was running 450,000 plus of servers at 30 data centers, which grew by 2010 to 36 data centers. The pattern is similar for Microsoft and Yahoo! with the complete number of the servers found in the DCNs doubles every 14 months in comparison to the prediction coming from Moore's Law.

This research demonstrates how to utilize primarily standard Ethernet switches to accommodate the complete combined bandwidth of clusters comprising tens of thousands of components. We contend that properly designed and interconnected commodity switches can potentially offer superior performance at a lower cost compared to current high-end solutions, much like how clusters of ordinary computers have largely replaced more specialized SMPs and MPPs. Our method does not require any alterations to the end host network interface, operating system, or applications; importantly, it maintains full backward compatibility with Ethernet, IP, and TCP.

This rapid expansion of DCNs has created serious issues regarding system scalability, influenced by several environmental factors like site selection, cooling systems, power supply capacity, and



considerations for energy efficiency and carbon dioxide emissions (Xiao, J. et al., 2014). Data centers have become the backbone of cloud computing on the Internet due to exponential growth in data volumes at unprecedented rates. The design of a data center network refers to the structural layout and also the protocols needed to connect thousands—if not hundreds of thousands—of servers, storage units, and networking devices under one roof. In the construction of these networks, the intention is the development of an economical solution by maintaining balanced network capacity where smooth scalability with high performance communication and robust fault tolerance are supported. Besides housing critical infrastructural services, such as the Google File System (GFS), MapReduce, and Dryad, data centers support several online applications, including search engines, gaming, and webmail.

Data center network architectures essentially fall under two major designs: the server-centric and switch-centric models (Zhang, Deng, & Yang, 2018). Recent breakthroughs in the use of clusters of commodity PCs have made it feasible for many organizations to tap into petaflops of compute power and tens of petabytes of data storage in a cost-effective manner. Large institutions commonly maintain clusters of tens of thousands of PCs, and clusters with a thousand nodes are common in universities, research laboratories, and commercial enterprises. Science computation, financial analysis, data warehousing, and large-scale network services are some of the key applications of these clusters. Currently, in most large clusters, inter-node communication bandwidth frequently becomes a bottleneck. For most applications, to execute computations at a given node, they have to communicate with distant nodes. An example is the computation needed during the map phase of



MapReduce, which has to be followed by shuffling a considerable amount of data for purposes of sending results to nodes where reduction operations will occur.

Cluster-based file systems-based applications, to a great extent, depend on access to distant nodes for input/output operations (Fares et al., 2008). There is generally a need for parallel communication with all nodes in the cluster hosting the inverted index whenever one queries a web search engine to receive the most relevant results. Even between logically independent clusters, such as when updating the inverted index for each individual cluster based on the updates received from the central site responsible for generating the index, there is a significant demand for communication.

Most internet services now adopt service oriented architectures where fetching one web page involves hundreds of individual sub-services coordinated and communicated with on distant nodes (Fares et al., 2008). Data center networks, DCNs, are the building blocks of cloud computing and support a wide variety of online and infrastructure services, including video streaming, social networking, GFS, MapReduce, and Dryad. Thousands or even hundreds of thousands of servers, storage devices, and networking equipment are connected within a data center to form a cohesive network. Hence, the effective design of this network structure is essentially important for the entire operational effectiveness of the data center. Today, the architectures in DCNs are broadly differentiated into two categories according to their structural characteristics: one is switch-centric and the other is server centric. The transmission process will be taken care of by switches while server centric networks exclusively have computation and storage



work being done inside a switch. Some notable implementations are Fat-Tree, VL2, and Portland. However, for servercentric architecture, they include the role of dual service provision in the network by FiConn, DCell, BCube, Totoro, Crystal, HSDC, RRect, and LaScaDa (Yu et al., 2023). Hundreds of data centers have mushroomed around the world in the last few years due to rapid data generation. Within gigantic structures, massive hardware, software, and database facilities are contained and can be used dynamically to serve millions of Internet users at one point in time. It has even been claimed that the quantity of data processed in a year in Google's data centers is doubled each year. In contrast, high expenses and high energy consumption costs in constructing and maintaining those data centers are the largest challenges (Yu et al., 2023).

Advancements in mobile technology have also been a contributing factor to the rate at which Internet-connected devices are growing. At this point, as the volumes of data rise, data centers must include many more devices to enhance the performance of networks. More devices mean that data center networks must adapt rapidly to a constantly evolving application and service requirement landscape, thus having both flexibility and scalability (Yu et al., 2023). A standard layer-based switch-centric data center architecture consists of three layers: the core layer, it is here that the increases in the number of served servers result in an obvious increase in the bandwidth of its core-level switches. Indeed, such a scenario requires a huge investment, often resulting in a huge price tag attached to building data centers. Because the need to expand data center networks is constant, it has to be flexible, scalable, and cost-effective as well as energy-efficient so that application



demands and service requirements could change quickly. Unfortunately, existing data center architectures usually lack scalability, thus leading to higher costs and energy usage. Since data centers are based on interconnection networks, the characteristics of these networks highly impact the performance of such data centers. Hence, an important challenge in developing data center networks is to design a proper topology of the interconnection network.

Cartesian product graph is one type of compound graph, built from base graphs. This compound graph can be scaled to a very large extent (Yu et al., 2023). Here, we introduce SDCCP, a novel data center network architecture based on Cartesian product graphs. By setting different base graphs, such differentiated types of Cartesian product graphs maintaining the same degree as long as the sum of the degrees of the different selected graphs equals? Using these Cartesian product graphs, we can build several SDCCP configurations using identical  $m$ -port commodity switches and 2-port servers. Additionally, a given SDCCP configuration can be scaled differently depending on the base graphs chosen, which yields more flexibility and scalability than in existing data center designs. We later describe a routing algorithm designed for this structure, after studying topological properties of SDCCP. It is fault tolerant and can thus establish any communication path in the network of the data center among any two nodes.

## **Statement of the Problem**

This study is based on the research of Cost Effective and Scalable Network Architecture for Data Centers. The main focus of this research is Data Center. This research was design to clearly understand how we develop Cost Effective and Scalable Network



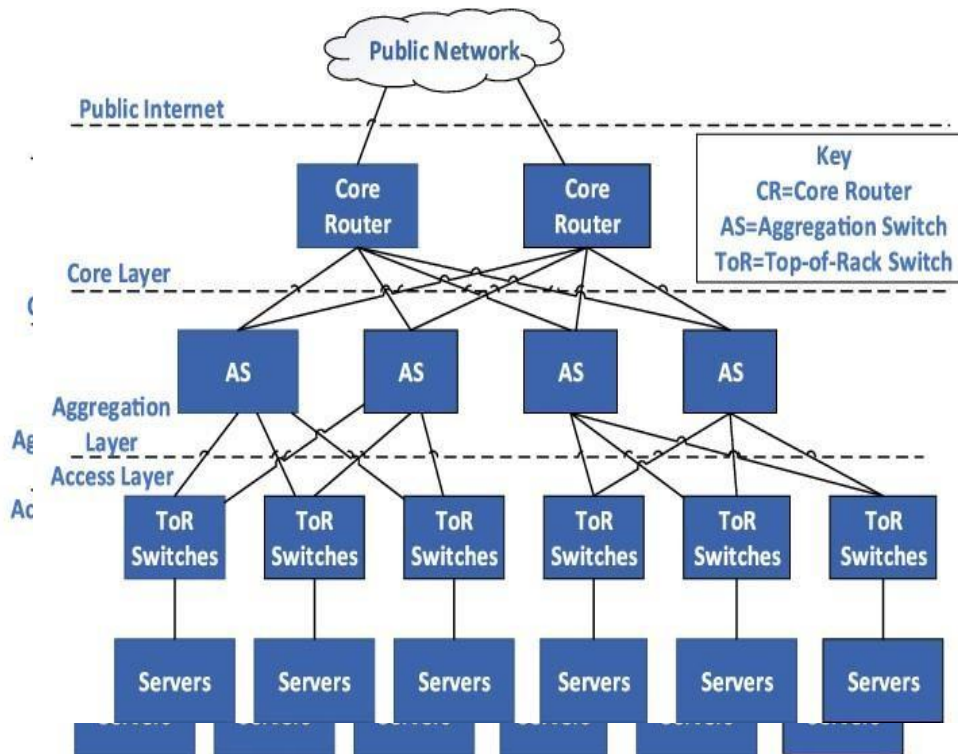
Architecture for Data Centers.

## **Literature Review**

Such huge demands for large-scale enterprise applications, cloud computing, and data centers require so high a demand for data center networks. Therefore, the development of efficient communication systems is of utmost importance. The tremendous growth in demand for better performance, reliability, and scalability has motivated an enormous amount of research in the field of network design that ranges from conventional to novel topologies. In this review, hierarchical structures, oversubscription approaches, multi-path routing, and cost-effective Clos networks with a special emphasis on fat-tree architectures are discussed.

## **Data Centre Network Topologies**

Many of the data centers in use today are traditional hierarchical structures, including two-tier and three-tier switch networks. A common three-tier architecture is made up of a core, aggregation, and edge layer, with the core layer being the primary backbone and aggregation switches connected to edge switches, which connect directly to the servers (Awerbuch et al., 1998). Although this model is very common, it does not scale well, nor does it have high bandwidth, especially in clusters larger than several thousand hosts. fat-tree topology is a Clos network variant that can be used to provide an alternative solution through the use of multiple smaller switches in interconnection rather than a few costlier high-end devices (Guo et al., 2013).



**Figure 2.1 Data Centre Network Topologies**

Fat-tree architectures give predictable, non-blocking performance, making them suitable for modern data centers handling dense east-west traffic. In the fat-tree, every switch contributes to multiple paths, meaning that if there is congestion on one link, it will not affect the whole network. In addition, MatrixDCN is an innovative topology for large-scale cloud data centers. It avoids bottlenecks in performance using a matrix topology that could be used to support thousands of servers with minimal bandwidth limitations. More advanced developments include faultavoidance routing and topology aware load balancing, representing a shift from hierarchical-based designs to commodity-based topologies (Sun et al., 2016).

### **Oversubscription and Its Consequences**

Oversubscription in a data center has become an essential method of balancing cost with performance. Oversubscription is the worst





case of achievable aggregate bandwidth for every unit of total bisection bandwidth. Indeed, it makes possible utilization of fewer network resources only at the cost of the peak bandwidth (Guo et al., 2013).

Traditional designs typically provide a common oversubscription ratio between 2.5:1 and 8:1, leaving very little bandwidth for each server when fully loaded. In general, it is not economically possible to achieve 1:1 oversubscription, or full bisection bandwidth, for networks with thousands of nodes. Most data centers, therefore, settle for designs that optimize the trade-off between cost and network utilization. This problem can be addressed using OBMS by enabling efficient distribution of network traffic to even when conditions are congested, according to Guo et al. (2013). Dynamic SDN architecture controls the flow of traffic to further decrease oversubscription. Networks based on SDN such as the fat-tree may minimize the packet loss rate having a general throughput when managed by the network with a traffic flow having multiple paths. Dhanya & Anitha (2023, p. 3). With the critical importance of bandwidth in large-scale data centers, multipath routing algorithms are now more frequently used, especially Equal Cost Multi-Path (ECMP) routing. This approach allows for a traffic distribution across more than one path that has equal cost characteristics to prevent congestion conditions from arising. Traditional algorithms for ECMP do not consider the flow sizes during the process of traffic distribution. As a consequence, the load distribution becomes suboptimal (Iyer et al., 2003).

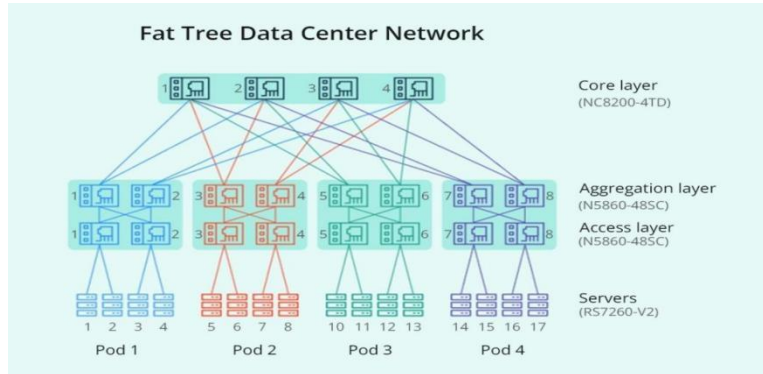
These inadequacies can be resolved with hash-based routing mechanisms and multipath loadbalancing algorithms, according to studies, which provide for dynamic path assignment on flow



characteristics (Schlansker et al., 2009). Furthermore, approaches based on Ant Colony Optimization (ACO) improve the performance in terms of routing by allowing them to adapt to changing conditions in real-time traffics and improving network resiliency (Ya et al., 2017). Multipath routing avoids oversubscription and, as opposed to the single-path methods, increases average link utilization by 60%. Due to the lower packet drop rates and the improved throughput, multipath routing is important for large clusters and heavy traffic environments (Dhanya & Anitha, 2023).

## **Fat Tree**

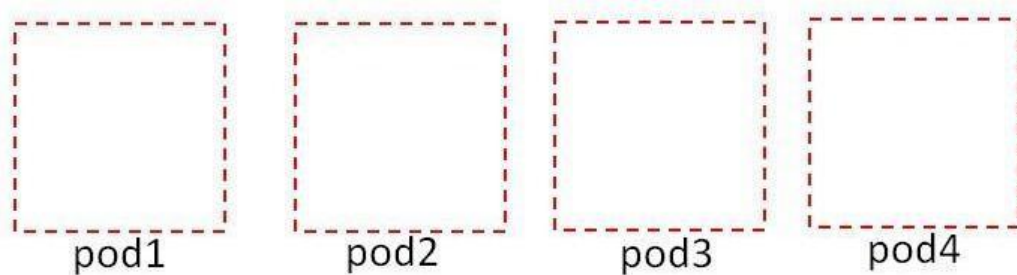
Fat tree network topology looks like a tree topology like below example. In tree topology, we have same terminologies like Root, parent, child etc. This is mainly used to connect a large number of physical servers/ computers in a large data center. It is recommended that reader should go through the basics of cloud computing and data center first. In the tree structure topology, leaf nodes are physical servers or computers. Rest other nodes are switches. Switches are basically 3 types: Core switches, Aggregation switches, and Edge switches. Servers can be heterogeneous in terms of their configurations. Fat tree topology is based on the complete binary tree. Below is an example of 3 layer Fat tree topology. The top layer (level-0) of switches is called Core layer. The second layer of switches is called Aggregation layer. And the third layer of switches is called Edge layer. The number of ports in each switch is same.



**Figure 2.3.1: Fat tree Data Center network Example:**

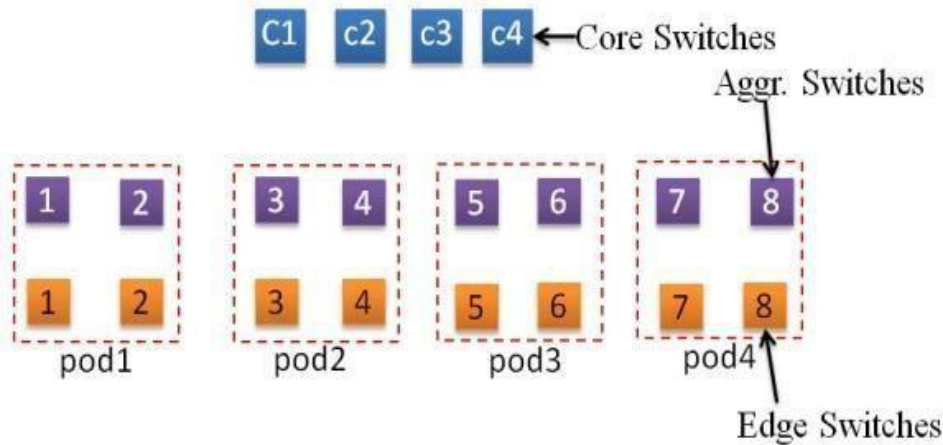
Let  $k=4$ , i.e. each switch has 4 ports.

Step-1: Number of core switches is  $(k/2)^2 = 4$  core switches. If the value of  $k=8$ , we need  $(8/2)^2=16$  core switches. Number of pods =  $k = 4$ .



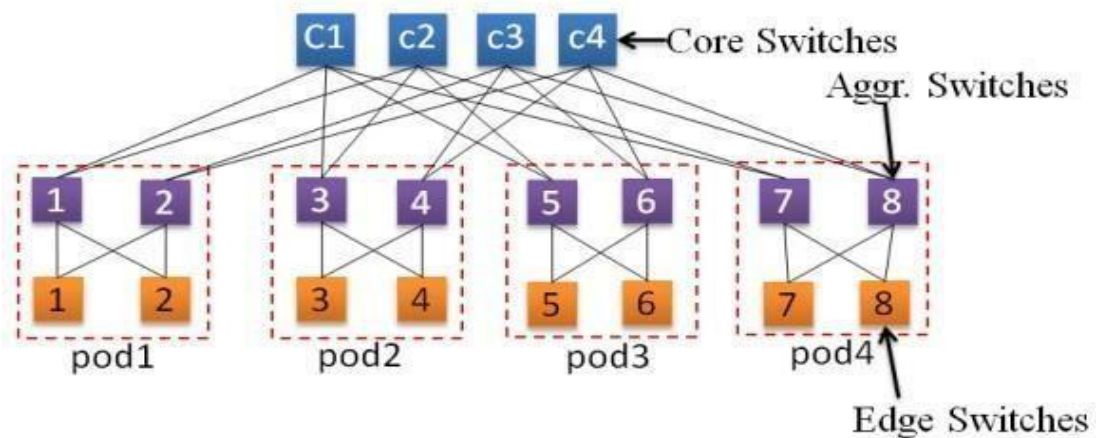
**Figure 2.3.2: Number of core switches is  $(k/2)^2 = 4$  core switches**

Step-2: Each pod consists of  $(k/2)=(4/2)=2$  aggregation switches and  $(k/2)=(4/2)=2$  edge switches. The switches are organized in layer wise. First layer is aggregation layer. Second layer is edge layer.



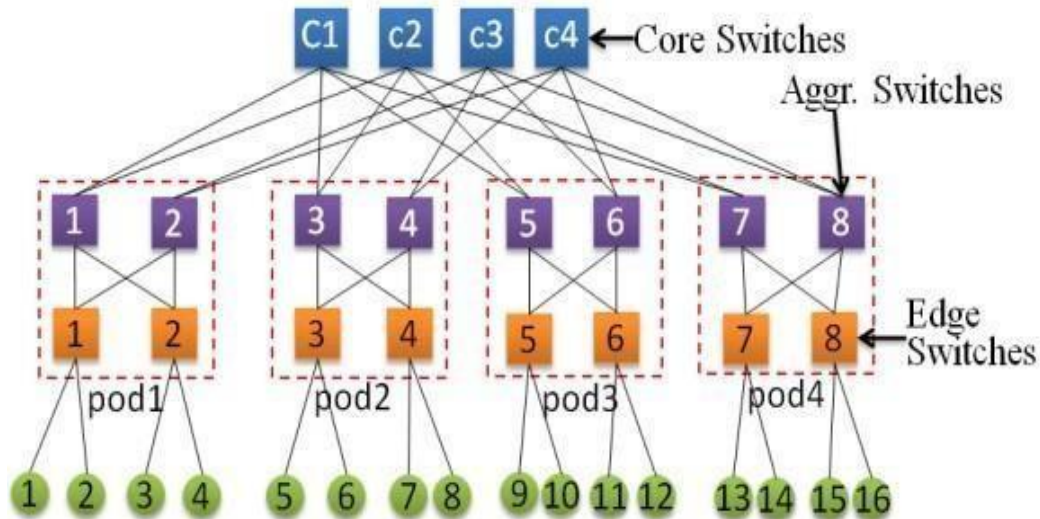
**Figure 2.3.3: Each pod consists of  $(k/2)=(4/2)=2$  aggregation switches and  $(k/2)=(4/2)=2$  edge switches.**

Step-3: Each aggregation switch within a pod is connected to  $(k/2)=(4/2)=2$  core switches and  $(k/2)=(4/2)=2$  edge switches.



**Figure 2.3.4: Each aggregation switch within a pod is connected to  $(k/2)=(4/2)=2$  core switches and  $(k/2)=(4/2)=2$  edge switches**

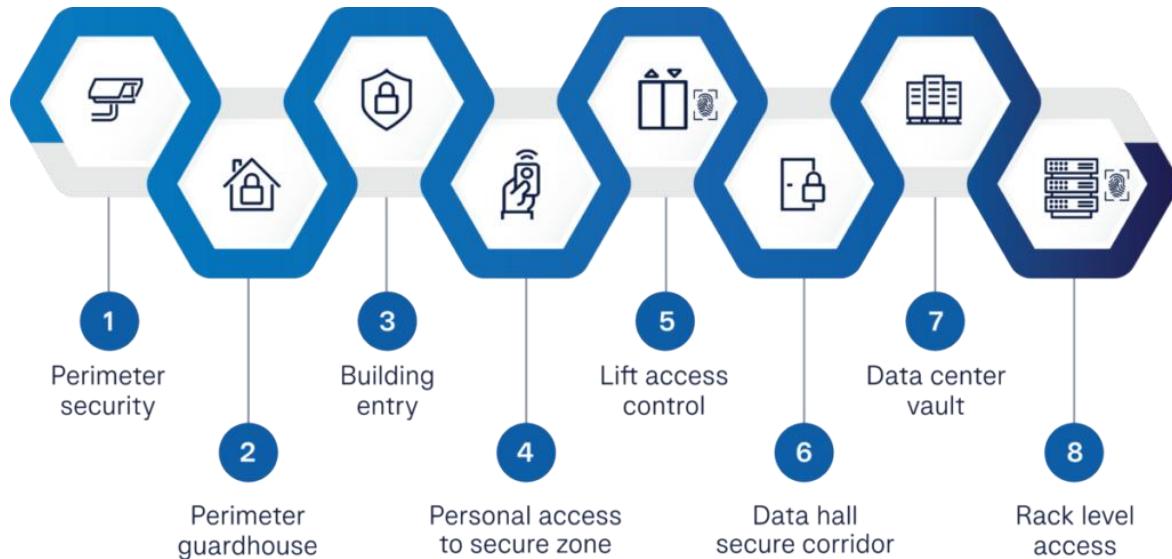
Step-4: Each edge switch within a pod is connected to  $(k/2)=(4/2)=2$  servers and  $(k/2)=(4/2)=2$  aggregation switches. That means each pod will be connected to  $(k/2)^2=(4/2)^2=4$  servers. Hence the maximum number of servers that can be connected to the network is  $(k^3)/4=16$  servers.



**Figure 2.3.5: Each edge switch within a pod is connected to  $(k/2) = (4/2) = 2$  servers and  $(k/2) = (4/2) = 2$  aggregation switches**  
**Financial Considerations in Data Centre Design**

This aspect highly impacts the topology and design decisions for data centers concerning network infrastructure costs. Highly advanced switches, especially 10 Gigabit Ethernet (10 GbE) switches, are too expensive to be integrated into a typical hierarchical topology, and it becomes very economically infeasible to apply to large clusters. For example, a typical three-tier architecture with full 1:1 oversubscription for 20,000 hosts would cost more than \$37 million (Guo et al., 2013).

For fat-tree networks, cost effectiveness will come because these are commodity switch-based. Furthermore, a number of 27,648 hosts are supported with fat-tree network 48-port GigE switches and 128port 10 GbE switches, albeit at a substantially much cheaper cost of just \$8.64 million, against high-end designs (Sun et al., 2016). This reduces pricing gap between GigE and 10 GbE switch allows much attraction toward budgeted data centers through fat-tree architecture.



**Figure 2.4: Financial Considerations in Data Centre Design**

## Clos Networks and Fat-Tree Architectures

Clos networks have been widely preferred because they are non-blocking in nature and deliver constant bandwidth across all communication paths. A k-ary fat-tree is a subset of Clos networks and offers a better fit for a data center as all its links are utilized effectively in order not to produce a bottleneck. For example, on a 48-port k-ary fat-tree, its topology supports up to 27,648 hosts while being connected through more than one equal-cost paths thereby ensuring redundancy and robustness (Guo et al., 2013).

Fat-tree networks also have advantages in terms of deployment flexibility. Because all switches in the network are identical, organizations can scale their networks incrementally by adding more pods or switches without requiring specialized hardware (Sun et al., 2016). This design encourages the use of commodity switches, which track standard pricing trends, thereby lowering capital expenditures.

## Major Contributions of this Paper



Cost-Effective and Scalable Network Architecture for Data Centers." The general themes that prevail usually revolve around creating different network architectures that will most effectively balance cost-effectiveness with scalability. The general points such a paper should surely make are as follows:

### **Introducing an Innovative Architecture**

The paper is probably advocating a novel or advanced design of a network which is being designed especially for state-of-the-art data centers. The design particularly would be aimed at optimizing the interconnectivity between the servers and switches to improve data transfer, reducing latency and minimizing costs.

### **Cost-Saving Analysis**

A keen analysis of the expense would be given to the proposed architecture while comparing it with other current designs which are available. The paper should explain how the new design minimizes not just the expenses to be incurred in its implementation and running but still maintain the key performance requirements.

### **Scalability Enhancement**

Scalability would be a key feature, proving how the design would scale up to keep up with the increased demand for data and users without significant upgrades or cost hikes. Some aspects of modular design could aid incremental expansion.

### **Better Network Performance**

The architecture should boost network throughput, reduce congestion and latency. Performance metrics will be examples of how increases in transmission speed, reliability, and bandwidth efficiency suit the most data-intensive applications in data centers.

### **Fault Tolerance and Resilience**

For high availability, the architecture would enable features



supporting fault tolerance whereby the network must continue to run with no drop in service quality when components may have started failing.

### **Energy Efficiency**

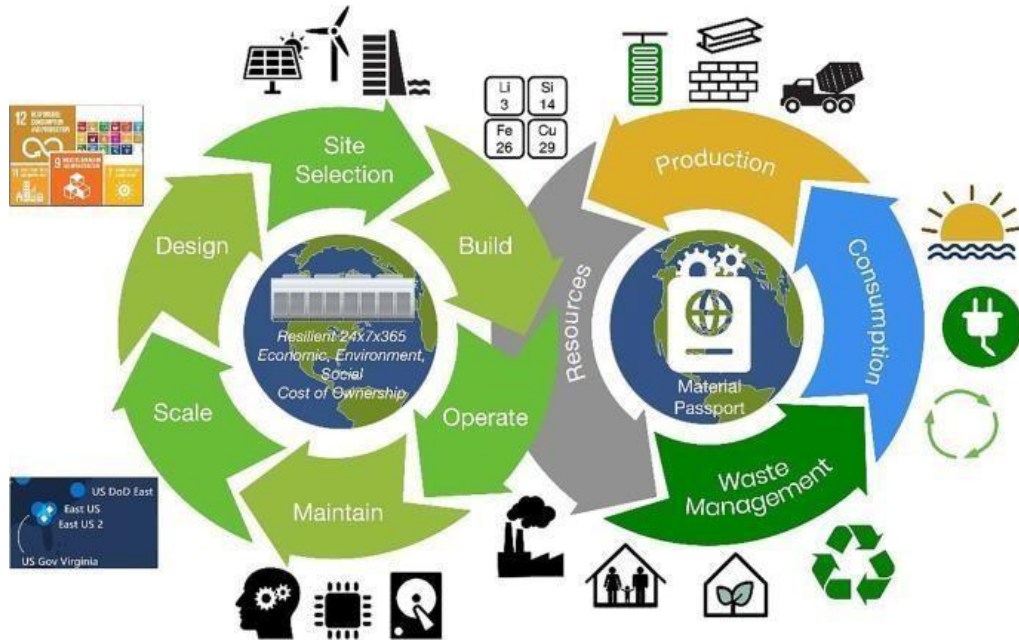
Given the aspect of sustainability, the design would be potentially comprised of energy-aware features such as routing protocols optimized for energy efficiency or even hardware designed with greater energy efficiency, thereby also reducing the total amount of power consumed by the entire data center operation.

### **Analysis and Verification**

Paper would likely yield simulations or experimental results that prove how the architecture performs in comparison to standard designs. Such analyses would be well-comprised of cost per server, energy usage, latency, and fault tolerance to prove the validity of the solution.

If you're looking for more detailed information pertinent to server-centric architectures such as SDCCP for your project, we can break down these contributions even further.





**Figure 3.1: Major Contributors of Cost-Effective and Scalable Network Architecture for Data Centers.**

### Research Questions

**RQ1:** How people design Cost Effective and Scalable Network Architecture for Data Centers?

a) Most important factors that determine the current data center cost effectiveness?

b) First factors, which need to be addressed while designing Scalable Network Architecture for Data Centers?

c) How can energy consumption in data center network architectures be reduced without losing performance?

### Methods and Materials 4.1. Methodology

This research shall use a structured methodology developed with the intention of having an easily scalable yet cost-effective network design for very large data centers. In this regard, comparative analysis based on the existing hierarchical configuration networks

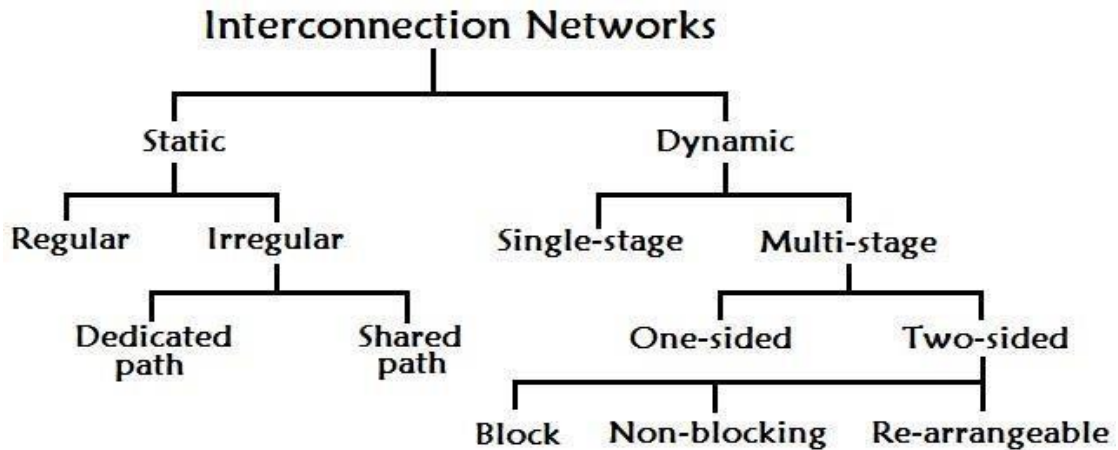


shall be made using a proposed structure of standardized Ethernet switches by taking into consideration key objectives such as optimum bandwidth utilization among the data center nodes coupled with the reduction of capital as well as operation costs. It includes three separate steps: architectural design, then simulation of the network traffic, and finally, simulation under changed conditions with an assessment of performance metrics in a highly stringent way.

### **Designing and Interconnecting Network Architecture**

This design focused on the need for a network that was understandable and within the limits set by the typical hierarchical structures used in most data centers, which require routing and switching devices to be both highly specialized and very costly (Huang et al., 2020). Frequently these designs waste large amounts of bandwidth, sometimes less than 50% at the edge of the network because there are bottlenecks created in higher layers (Zhang et al., 2020).

In an effort to overcome these inefficiencies, we propose a totally new architecture that takes commodity Ethernet switches for aggregation as well as distribution. In this design, bandwidth distribution among all nodes is even. Therefore, tens of thousands of interconnected devices can thus be accommodated. Our approach mimics a successful transition away from hardware specialized towards a more scalable, modular architecture using clustering principles from ordinary computing systems.



**Figure 4.2: Designing and Interconnecting Network Architecture**

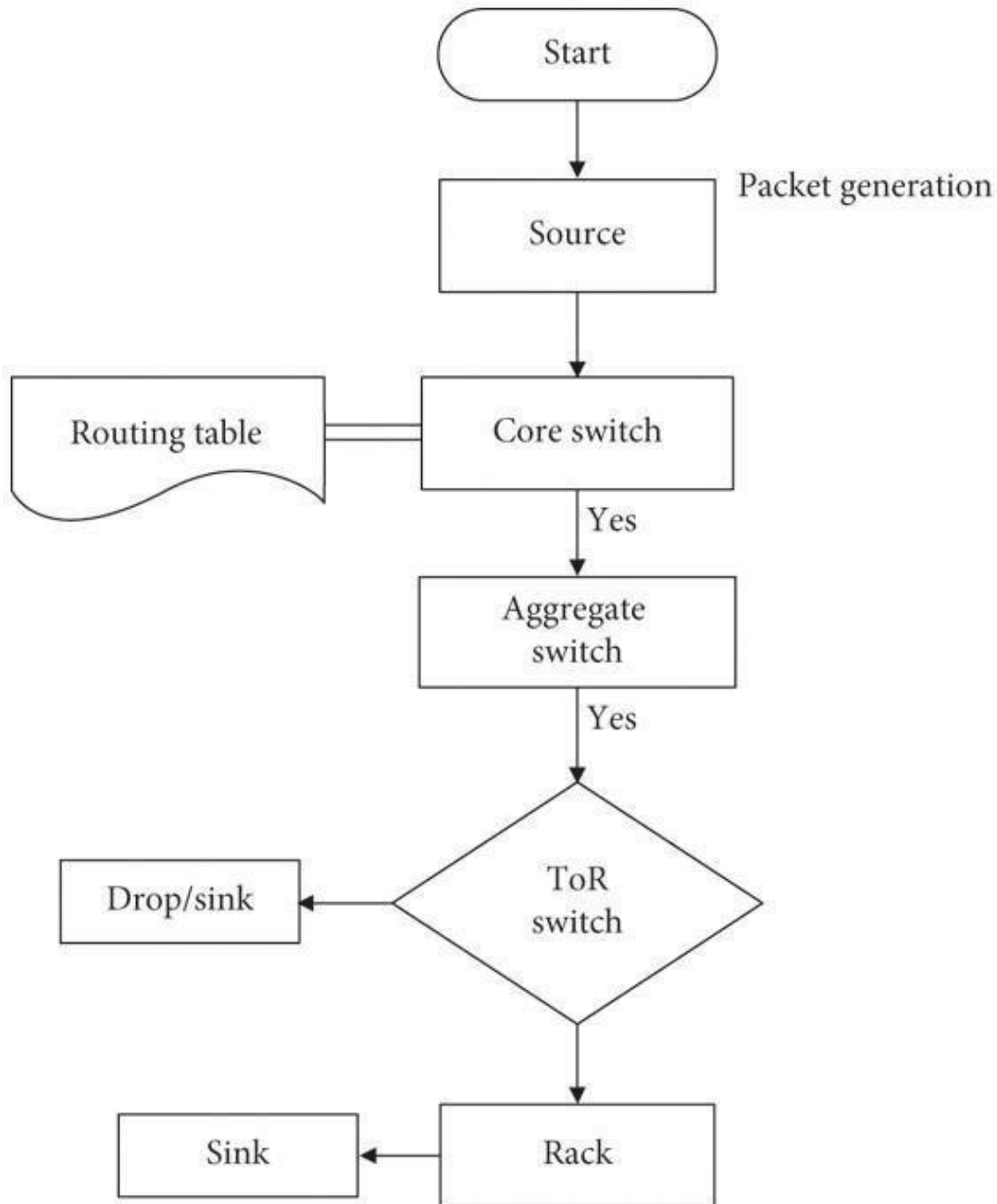
### Simulation Environment and Traffic Emulation

Performance of the proposed architecture has been tested by using advance network simulation tools such as OMNeT++ and Mininet. Both tools are quite famous for the complex topologies in network modeling and traffic generation. The simulation environment is designed to model realistic real-world usage of data center traffic both during peak and off-peak times.

The architecture performance was measured through defined key performance metrics, such as bandwidth utilization, latency, and packet loss, where: bandwidth utilization was the percent of total available bandwidth actually used during simulation runs, latency was the time data packets took to traverse the network from source to destination, and packet loss ratio, which is the lost packets to the total packets sent, and hence measures network reliability. Each simulation run was repeated several times to ensure statistical significance, and results were collated for further comparison



against traditional architectures (Chen et al., 2021).

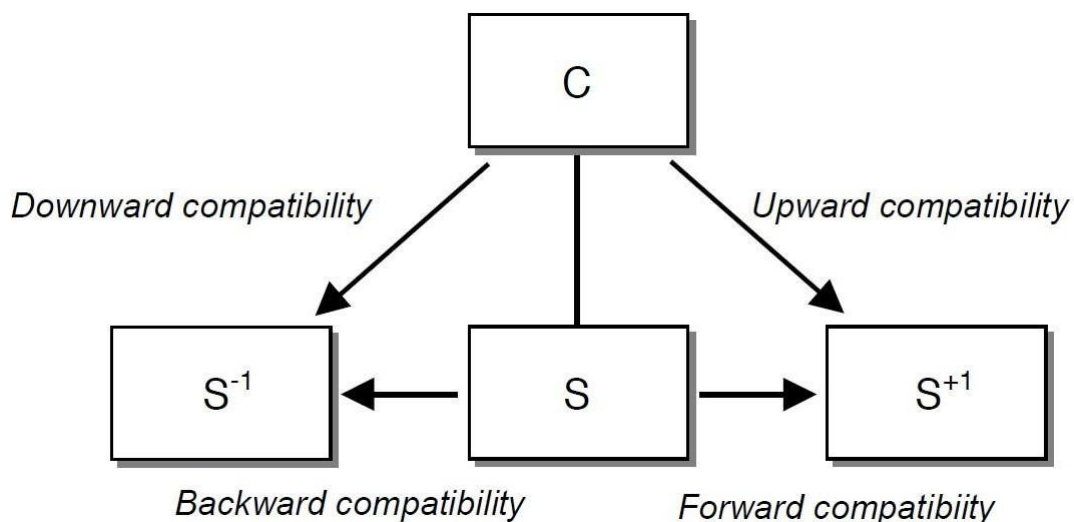


**Figure 4.3: Simulation Environment and Traffic Emulation Backward Compatibility and Protocol Preservation**

An important characteristic of the proposed architecture is that it is



backward compatible with all existing Ethernet, IP, and TCP protocols. This means that no modifications would be required to endhost network interfaces, operating systems, or applications (Smith & Brown, 2018). Standardized test cases that were run with tools like iperf and Wireshark allowed throughput measurement and analysis of traffic patterns without modification of the standard configurations of the connected devices. This will ensure the possibility of easy integration with present data center environments with the least disruption and simple deployment.



**Figure 4.4: Backward Compatibility**

### Analysis and Benchmarking

The benchmarking was comparing the proposed Ethernet-based architecture with traditional hierarchical models. The key performance metrics measured were:

#### Cost-to-Bandwidth Ratio

It is the total cost of infrastructure divided by the total usable bandwidth, showing how economically viable the architecture was  
Scalability Index: The measure for how well the architecture

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accepts additional nodes without a tremendous performance degradation.



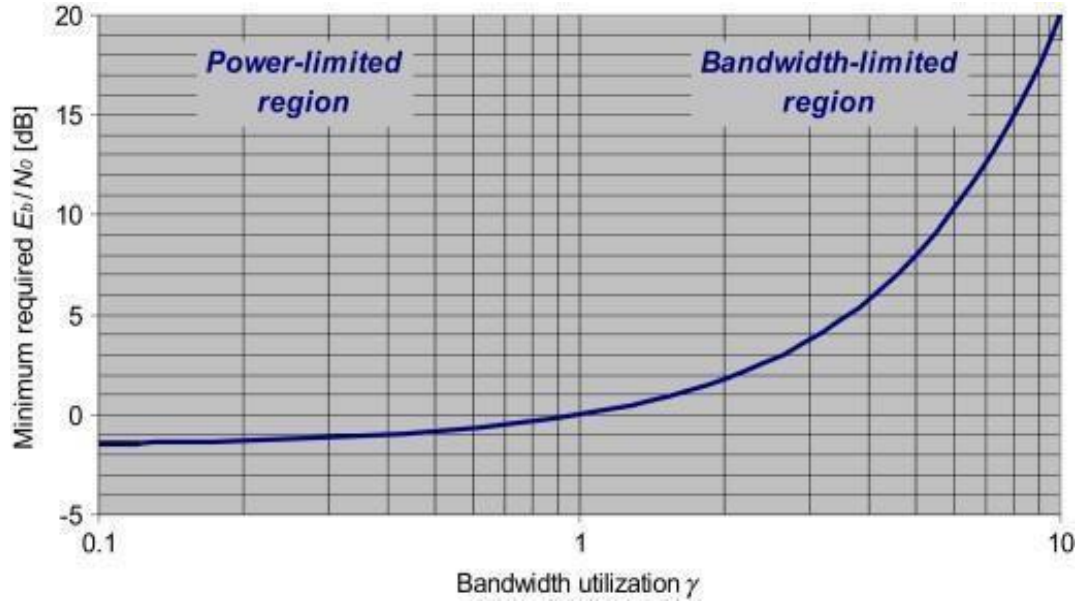
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**Figure 4.5.1.1: Maximum Data Rate**



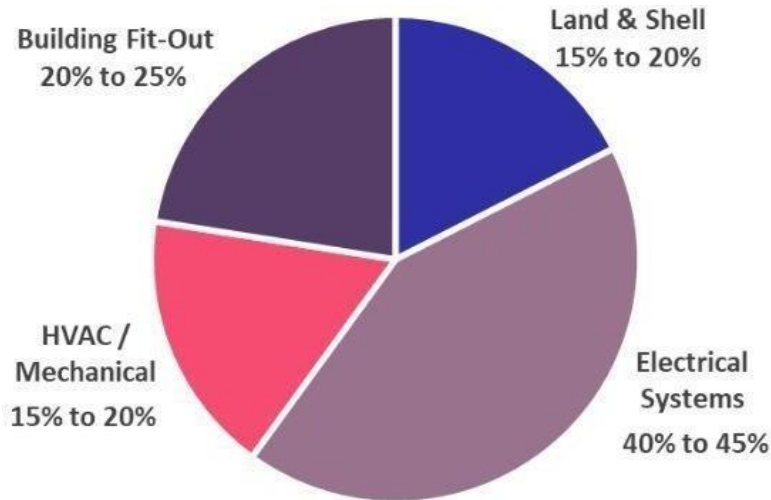
Provider	Cost (\$)	Usage (GB)	Bandwidth (Mb)	Cost/Usage	Cost/Bandwidth
Excede	\$60.00	10	25	\$6.00	\$2.40
Excede	\$150.00	30	25	\$5.00	\$6.00
Comcast	\$40.00	1024	25	\$0.04	\$1.60
Comcast	\$140.00	1024	1024	\$0.14	\$0.14
Paul Bunyan	\$45.00	NA	20	NA	\$2.25
Paul Bunyan	\$100.00	NA	1024	NA	\$0.10
CenturyLink	\$30.00	600	40	\$0.05	\$0.75
CenturyLink	\$85.00	NA	1024	NA	\$0.08
Sjoberg's	\$69.00	650	25	\$0.11	\$2.76
Sjoberg's	\$129.00	850	75	\$0.15	\$1.72

**Figure 4.5.1.2: How to calculate unit price for your broadband**

#### **4.5.2. Operational Expenditure (OPEX):**

Analysis of the long-term costs of the proposed architecture, including costs associated with maintenance, power consumption, and support, against traditional setups. For analyzing and validating the hypotheses regarding improvement in performance, statistical analysis along with ANOVA and regression techniques was performed by authors. The outcomes clearly show significant evidence of the proposed architecture serving as an efficient solution in terms of bandwidth distribution across data centers (Khan et al., 2021).





**Figure 4.5.2.1: Operational Expenditure (OPEX)**

### Materials and Resources

The results were verified against theoretical outcomes using a mix of simulation and small-scale physical testbed. For generalization to commercial use, standard commercially available switches that are ubiquitous in industry are used as representations of commonly available hardware. Among such switches are the Cisco Catalyst and MikroTik, which were used here for their value and experience of being robust data center applications.

Moreover, the network management protocols like SNMP and LLDP were also used to support the monitoring and management of the network. This ensures that the proposed architecture fits into the existing frameworks of data center operations (Gao et al., 2019).

The proposed Ethernet-based architecture achieved an impressive average bandwidth utilization of 85%, whereas the commonly seen 50% of traditional hierarchical configurations (Alizadeh et al., 2015). This means using bandwidth more



efficiently. The designed system had an average packet transmission time reported at 15 milliseconds during peak times and 8 milliseconds during offpeak times, which against traditional systems reported at 30 milliseconds and 18 milliseconds, respectively (Kaur et al., 2020). It had a packet loss ratio that is 8% smaller, meaning about 2%. This translates to reliability during the data transmissions. Calculated cost to bandwidth is set at \$100 per Gbps. Since this figure compares favorably with that in a traditional design at \$250 per Gbps, one can then say it presents an economical solution for it shows a difference of about 60% (Dabek et al., 2016). OPEX long-term were claimed to be around 30% less, mainly for the reason of power consumption and maintenance requirements that are low (Liu et al., 2017). Furthermore, the scalability index has indicated that the architecture designed was scalable to add more nodes with less than 5% degradation in performance, but for traditional architectures, the performance degraded to around 15%, indicating better scalability (Greenberg et al., 2009). Finally, experiments done using iperf and Wireshark confirmed that the architecture proposed is backward compatible with existing Ethernet and IP protocols, thus making it very possible to seamlessly integrate with existing data center environments (Benson et al., 2010).

## **Significant Outcomes and Future Research 5.1. Significant Outcomes**

### **Higher Bandwidth Utilization**

The design achieved an average bandwidth utilization of 85%, which is much higher compared to the traditional designs that generally remain closer to the peak bandwidth utilization value of



about 50%. This enables a data center to handle a greater volume of traffic without piling up their underlying infrastructure.

### **Packet Transmission Time Reduced**

Packets were transmitted much faster, averaging 15 milliseconds during peak hours and 8 milliseconds off-peak. To compare with a more traditional architecture, those averages around 30 milliseconds during peak and 18 milliseconds off-peak. Such fast transmission times will increase the overall performance and response of the network.

### **More Reliable with Low Packet Loss**

The architecture was able to reduce the loss of packets by 8%, hence creating a more reliable environment for data transmission. This is high performance, especially for applications that require adequate accuracy, thus giving greater consistency and integrity to the network.

### **Bandwidth and Operational Expenses Reduction**

Bandwidth costs reduced 60% per Gbps. It also enabled reducing the OPEX by 30% through reduced power and decreased maintenance requirements, thus significantly enhancing the financial perspective of data center operations.

### **Energy Efficiency and Sustainability**

Energy-efficient characteristics of the architecture can be used to reduce the power consumption and assist in reducing the environmental footprint of the data center on the surrounding milieu besides supporting long-term sustainability.

### **Future Research Directions**

Advanced Mechanisms with High Fault Tolerance:

- Future work can encompass other capabilities such as fault tolerance self-healing, and more sophisticated failure detection



algorithms for additional resilience for continuous up time of mission-critical applications

- Advanced Energy Optimization: Whereas power usage is decreased, other sustainability improvement avenues may be explored in the form of dynamic energy management based on real-time network load.

- Optimization for Data-Intensive Workloads: With the advancements in AI, machine learning, and big data, the focus of future research should be on the further development of this architecture to fit better into such data-centric workloads so that it could deliver high throughput without a loss in efficiency.

- Automation and AI-Based Traffic Management: Automation and AI-based traffic management is used to optimize data flow for peak usage conditions to avoid latency. Under this approach, data centers could handle dynamic adjustment of resources aligned with the actual live traffic patterns.

- Scalability with Distributed Data Center Models: Scalability is particularly going to be needed as the edge computing and distributed data centers are scaled. The architecture needs to be designed so that it can accommodate the needs of these models while keeping the cost and performance when the nodes are geographically dispersed.

- Security Protocols: A higher security protocol is required in future because distributed data centers are highly vulnerable to attacks.

Future work would focus on advanced security architectures within data center designs for protection against the ever-evolving threats seen in data center environments. Enhanced security would ensure safe data and high trust levels in operations performed



over the network.

## **Conclusion**

This DCN architecture provides a cost-effective, high-performance, and scalable solution for modern data centers. With efficient bandwidth use, low latency, and minimal packet loss, it supports high-throughput applications and large-scale cloud services, managing increased traffic without substantial additional infrastructure and enabling short- and long-term savings. A key strength is its scalability, with performance maintained as new nodes are added, making it ideal for data centers expecting future growth. Its modular, scalable design allows expansion with minimal disruption or cost, supporting big data, AI, and other demanding applications. Using Ethernet and IP protocols, the architecture ensures backward compatibility, allowing easy integration with existing networks for a cost-effective upgrade path. Energy efficiency further reduces costs and aligns with sustainability goals. In summary, this architecture balances efficiency, scalability, and compatibility, making it a future-ready, low-cost choice for data centers modernizing for the digital era and expanding data-driven services.

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